



Sheet 1

Problem 1.1 *Properties of almost every graph.*

For two positive integers a, b we say that a graph $G = (V, E)$ is (a, b) -extendable if the following holds: for all sets $A, B \subset V$ with $|A| = a$, $|B| = b$ and $A \cap B = \emptyset$ there exists a vertex $x \in V \setminus (A \cup B)$ such that x is adjacent to all vertices in A and to none in B .

- (1) Prove: for every fixed $(a, b) \in \mathbb{N}^2$, almost every graph is (a, b) -extendable.
- (2) Let $k \geq 1$ be an arbitrary fixed integer. Prove or disprove:
 - (i) Almost every graph is not k -connected.
 - (ii) Almost every graph has minimum degree at least k .
 - (iii) Almost every graph contains a separating clique.

Problem 1.2 *Large cuts in directed graphs.*

Let G denote an oriented graph (i.e. an ordinary graph together with one of two choices of an orientation for each edge). Prove that there exists a partition $A \cup B$ of $V(G)$ such that the number of edges from A to B is at least $\frac{1}{4} \|G\|$.

Problem 1.3 *Square of the expectation, expectation of the square, and the second-moment method.*

Let A_n be a sequence of discrete random variables with non-negative values with respect to some probability space \mathcal{P} . Denote by Ex and Var the expectation and variance w.r.t. \mathcal{P} . Prove or disprove: $\Pr[A_n \geq 1] \xrightarrow{n \rightarrow \infty} 1 \iff \frac{\text{Var}[A_n]}{\text{Ex}[A_n]^2} \xrightarrow{n \rightarrow \infty} 0 \iff \frac{\text{Var}[A_n^2]}{\text{Ex}[A_n^2]^2} \xrightarrow{n \rightarrow \infty} 0$

Problem 1.4 *Towards the clique number of almost every graph.*

For every $K \subset [n]$ let $\mathbf{1}_{n,K}: \mathcal{G}_n \rightarrow \{0, 1\}$ be defined by $\mathbf{1}_K(G) = 1 \iff G[K]$ is a complete graph. Let $A_{n,K} := \{G \in \mathcal{G}_n : G[K] \text{ is a complete graph}\}$. Let $\Pr_{\frac{1}{2}}$ denote the measure of $G(n, \frac{1}{2})$.

- (1) with $A_{s,k,n} := A_{n,[k]} \cap A_{n,\{k-s+1, k-s+2, \dots, 2k-s\}}$, which of the three following equations is true for every $1 \leq k \leq n$:

$$(i) \quad \sum_{\{C_1, C_2\} \in \binom{[n]}{2}} \Pr_{\frac{1}{2}}[A_{n,C_1} \cap A_{n,C_2}] = \sum_{0 \leq s \leq k-1} \binom{n}{s} \binom{n-s}{k-s} \binom{n-k}{k-s} \Pr_{\frac{1}{2}}[A_{s,k,n}] \quad ?$$

$$(ii) \quad \sum_{\{C_1, C_2\} \in \binom{[n]}{2}} \Pr_{\frac{1}{2}}[A_{n,C_1} \cap A_{n,C_2}] = \sum_{0 \leq s \leq k-1} \binom{n}{s} \binom{n-s}{\frac{k-s}{2}} \Pr_{\frac{1}{2}}[A_{s,k,n}] \quad ?$$

$$(iii) \quad \sum_{\{C_1, C_2\} \in \binom{[n]}{2}} \Pr_{\frac{1}{2}}[A_{n, C_1} \cap A_{n, C_2}] = \sum_{0 \leq s \leq k-1} \frac{1}{2} \binom{n}{s} \binom{n-s}{k-s} \binom{n-k}{k-s} \Pr_{\frac{1}{2}}[A_{s, k, n}] \quad ?$$

(2) which of the three following formulas equals $\Pr_{\frac{1}{2}}[A_{n, [k]} \cap A_{n, \{k-s+1, k-s+2, \dots, 2k-s\}}]$:
 $(\frac{1}{2})^{\binom{k}{2} - \binom{s}{2}}$ or $(\frac{1}{2})^{2\binom{k}{2} - \binom{s}{2}}$ or $(\frac{1}{2})^{\binom{2k-s}{2}}$?

(3) Prove or disprove: $\binom{n}{s} \binom{n-s}{k-s} \binom{n-k}{k-s} = \binom{n}{k} \binom{k}{s} \binom{n-k}{k-s}$ for every $1 \leq k \leq n$ and every $1 \leq s \leq k$.
 If you prove this, then your proof must be a ‘bijective’ one (not by simplifying formulas), i.e. proceed by describing two finite sets in bijective correspondence such that the left-hand side obviously counts one of them and the right-hand side the other.

(4) Let $\omega(\cdot)$ denote the clique number. Prove that $\omega(G) \geq \log|G|$ for almost every graph G .

Reminder. $[n] := \{1, \dots, n\}$, $|G| :=$ number of vertices of G , $\|G\| :=$ number of edges of G , $\mathcal{G}_n :=$ set of all graphs on $[n]$