

Due Monday, 26 Oct. 2015.

1: Let $H_{u,v}$ denote the number of complex roots of $e^{\sqrt{2}y} = 1$ in the horizontal strip $\mathbb{R} \times [u, v]$. Please find a positive integer C with $|H_{u,v} - \sqrt{2} \cdot \frac{v-u}{2\pi}| \leq C$ for all $u < v$.

Solution

First we find the solutions of $e^{\sqrt{2}y} = 1$, they are $\{\sqrt{2}k\pi i \mid k \in \mathbb{Z}\}$, all placed on the line $\operatorname{Re}(y) = 0$. So in every horizontal strip $\mathbb{R} \times [u, v]$ we have at most $\frac{v-u}{\sqrt{2}\pi} + 1$ solutions. Hence, we can choose $C = 1$.

2: Use induction, and Rolle's Theorem, to prove the following upper bound:

Any univariate polynomial f of the form $c_1x^{a_1} + \dots + c_t x^{a_t}$, with $c_i \in \mathbb{R}^*$ for all i and $a_1 < \dots < a_t$, has at most $t - 1$ positive roots.

Can you find, for any t , a univariate t -nomial with exactly $t - 1$ (distinct) positive roots?

Solution

We prove the statement by induction on t .

For $t = 1$, $f(x) = c_1x^{a_1}$, so its only root is 0. For $t \geq 2$, we can assume $a_1 = 0$, as otherwise we could factor $f(x) = x^{a_1}(c_1 + c_2x^{a_2-a_1}, \dots, c_t x^{a_t-a_1})$ and the factor x^{a_1} could be left out, as it does not contribute to the number of positive roots of f . Assume f has N distinct positive roots. They define $N - 1$ open intervals where we can apply Rolle's theorem, and obtain that f' has at least $N - 1$ positive roots. As f' has $t - 1$ terms, by induction it must be $N - 1 \leq t - 2$, which means $N \leq t - 1$.

As t -nomials with exactly $t - 1$ positive roots, choose $f(x) := (x - 1)(x - 2) \cdots (x - t + 1)$.

3: Suppose $f(x) = 3^{\beta+\gamma}a - 3^\beta b x^\alpha + c x^{\alpha+\beta}$ for some $a, b, c, \alpha, \beta, \gamma \in \mathbb{N}$ with $\gamma > \alpha$ and $3 \nmid a, b, c$. Please prove that if f has an integer root of the form $3^r s$ for some $r \in \mathbb{N}$ and $s \in \mathbb{Z}$, with s not divisible by 3, then there are just two possibilities for r . Please find these two possible values for r .

Solution

If $f(3^r s) = 3^{\beta+\gamma}a - 3^\beta b(3^r s)^\alpha + c(3^r s)^{\alpha+\beta} = 0$, then $3^{\beta+\gamma}a + 3^{\alpha r + \beta r} c s^{\alpha+\beta} = 3^{\alpha r + \beta} b s^\alpha$. Both sides must contain the same power of 3 in their factorization, and as $3 \nmid a, b, c, s$ it follows

$$\min\{\beta + \gamma, \alpha r + \beta r\} = \alpha r + \beta$$

If $\min\{\beta + \gamma, \alpha r + \beta r\} = \beta + \gamma$, then $r = \frac{\gamma}{\alpha}$, if $\min\{\beta + \gamma, \alpha r + \beta r\} = \alpha r + \beta r$, then $r = 1$.

4: Please prove that, given any *finite* subset $X \subset \mathbb{R}^2$, the set $\mathbb{R}^2 \setminus X$ is path-connected. In particular, please outline your own method to build an explicit path in $\mathbb{R}^2 \setminus X$ connecting any two given points p and q . You should think of your construction as an algorithm you could actually implement.

Solution

Suppose $p = (a, b), q = (c, d) \in \mathbb{R}^2 \setminus X$, define a family of parabolas \mathcal{P}_α passing through p, q , that if $a \neq c$ are defined by the equations

$$y = \alpha x^2 + \beta x + \gamma \quad \alpha \in \mathbb{R}$$

where

$$\beta = \frac{d-b}{c-a} - \alpha(c+a) \quad \gamma = b - \alpha a^2 - a \frac{d-b - \alpha(c^2 - a^2)}{c-a}$$

(if $a = c$ use the family $x = \alpha y^2 + \beta y + \gamma$, and suitable β, γ instead). One can choose α so that \mathcal{P}_α does not pass through any point of X , and use the arc of parabola (p, q) as path.

- 5:** (a) Please prove that $\log |x^A| = (\log |x|)A$ for any $x = (x_1, \dots, x_n) \in (\mathbb{C}^*)^n$.
 (b) Please prove that $f(x) := x^A$ is invertible (as a function from \mathbb{R}_+^n to \mathbb{R}_+^n) $\iff \det A \neq 0$.

Solution

- (a) Following the definition given in the lecture, denoting c_1, \dots, c_n the columns of $A \in \mathbb{R}^{n \times n}$, we have

$$\begin{aligned} \log |x^A| &= \log |(x^{c_1}, \dots, x^{c_n})| = (\log |x^{c_1}|, \dots, \log |x^{c_n}|) = \\ &= (\log |x_1^{a_{1,1}}| + \dots + \log |x_n^{a_{n,1}}|, \dots, \log |x_1^{a_{1,n}}| + \dots + \log |x_n^{a_{n,n}}|) = \\ &= (a_{1,1} \log |x_1| + \dots + a_{n,1} \log |x_n|, \dots, a_{1,n} \log |x_1| + \dots + a_{n,n} \log |x_n|) = (\log |x|)A \end{aligned}$$

- (b) " \Leftarrow " If $\det(A) \neq 0$, then define $f^{-1}(x) = x^{A^{-1}}$, that is the inverse of f by the proposition stated in class.
 " \Rightarrow " We show that if $\det(A) = 0$, then f is not invertible. Since A is not invertible, it has a nonzero left null vector $v = (v_1, \dots, v_n)$. In particular, $f((1, \dots, 1)) = 1 = f(2^v) = f((2^{v_1}, \dots, 2^{v_n})) = 2^{vA} = 2^{(0, \dots, 0)}$. Since $(1, \dots, 1)$ and 2^v are distinct, we are done.

NOTE: Please feel free to e-mail comments, questions, and/or corrections.