

Lecture: Scheduling

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Organisational issues

- Lectures: Wed 8:30 – 10:00 in Room 00.06.011 (MI Hörsaal 3)
- Exercises (by Roman Rischke)
 - Wed 10:15 – 11:45 in Room 02.04.011 (seminar room)
 - begin next week
 - two groups in alternating weeks
- Lecture notes: electronic blackboard
- Homepage: <http://www-m9.ma.tum.de/WS2015/Scheduling>
- Exam: oral/written (will be decided later)

Content of this course

- Classification of scheduling problems
- Focus on standard deterministic scheduling problems
- Complexity of scheduling problems
- State of the art solution methods and analysis of algorithms
- Exact and approximation algorithms

Scheduling models are motivated by real-world applications. In this course we focus on theoretical foundations.

Literature

books

- M.L. Pinedo (2012). *Scheduling. Theory, Algorithms, and Systems*. Springer.
- P. Brucker (2007). *Scheduling Algorithms*. Springer.
- F. Jaehn and E. Pesch (2014), *Ablaufplanung. Einführung in Scheduling*. Springer.

articles

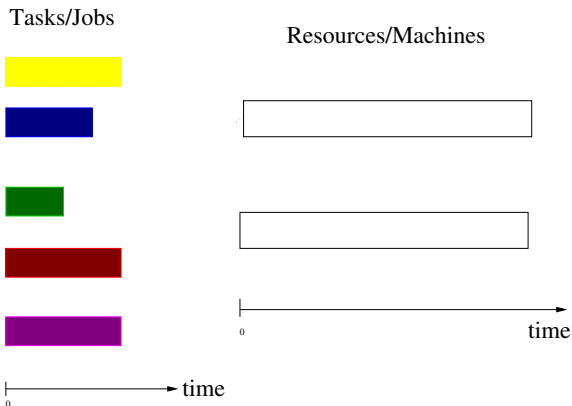
- R.L. Graham, E.L. Lawler, J.K. Lenstra, A.H.G. Rinnooy Kan (1979). Optimization and approximation in deterministic sequencing and scheduling: a survey. *Annals of Discrete Mathematics* **5**, pp. 287–326.
- E.L. Lawler, J.K. Lenstra, A.H.G. Rinnooy Kan, D.B. Shmoys (1993). Sequencing and scheduling: Algorithms and complexity. Chapter 9 in S.C. Graves, A.H.G. Rinnooy Kan, P.H. Zipkin. *Logistics of production and inventory*. Handbooks in Operations Research and Management Science **4**.
- B. Chen, C.N. Potts, G.J. Woeginger (1998). A review of machine scheduling: complexity, algorithms and approximability. In D.Z. Du and P.M. Pardalos, *Handbook of Combinatorial Optimization*, pp. 21–169.

Introduction and Overview

What is scheduling?

Definition (Lawler, Lenstra, Rinnooy Kan and Shmoys (1993))

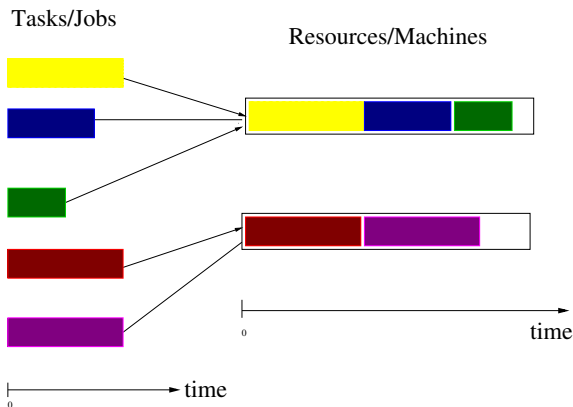
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Elements

Resources:

Machines, staff, crane, construction crews, CPUs, ...

Jobs:

Operations in a production process, stages in a construction project, execution of computer programs, ...

Objectives:

- Minimizing last completion time
- Maximize number of tasks completed on time
- ...

Example 1: Gate assignment at an airport

- Airport: dozens of gates, hundreds of planes arriving/departing per day
- Flight schedule: specifies arrival/departure time, passenger demand
- Service time at gate: disembarking, service of the plane, boarding
- Task: assign planes to gates while minimizing objectives such as work for airline personnel and airplane delays



Example 2: Staff scheduling in hospitals

- In a hospital, staff is working night and weekend shifts
- Due to labor agreements and law, not too many consecutive night shifts can be done by one staff member
- Shifts need to be divided over staff in a fair manner
- Staff members can submit their preferences
- First goal: feasible schedule
- Second goal: fair schedule

Example 3: Scheduling tasks in a CPU

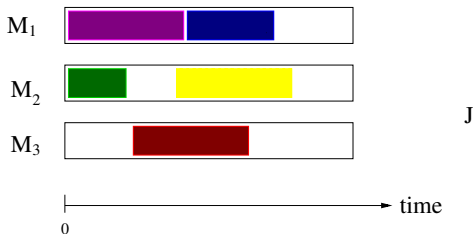
- CPU = central processing unit
- Allocate CPU time to tasks.
- Goal: Minimize average time a tasks spends in the system.



General scheduling models – representation

- m machines, M_1, \dots, M_m or $1, \dots, m$
- n jobs, J_1, \dots, J_n or $1, \dots, n$
- A job j has processing time p_{ij} when processing on machine i .
- Schedules represented by **Gantt charts** (developed by Henry Gantt in the 1910s to illustrate project schedules.)

Gantt chart for 3 machines, 5 jobs



Classification of scheduling problems

Classification of scheduling problems

Many scheduling problems can be described by a three-field notation

$\alpha | \beta | \gamma$:

α denotes machine environment

β denotes job characteristics

γ denotes objective function

Remark

A field may contain more than one entry, but may also be empty.

Classification – machine environment (α)

- $\alpha = 1$: single machine
- $\alpha = P$: identical parallel machines
 - m identical machines
 - processing time of job is independent of machine: $p_{ij} = p_j$
- $\alpha = Q$: uniform or related parallel machines
 - m parallel machines with speeds s_1, \dots, s_m
 - processing time $p_{ij} = p_j/s_i$
 - if $s_1 = \dots = s_m$ then case of identical machines
- $\alpha = R$: unrelated parallel machines
 - m different parallel machines
 - processing times arbitrary p_{ij}
- $\alpha = Pm, Qm, Rm$: number of machines is m (fixed)
- ...

Classification – job characteristics (β)

β may contain several values

- r_j : release dates
 - if $r_j \in \beta$ then job may not start before time r_j
 - if $r_j \notin \beta$ then job may start at any time
- $pmtn$: preemption
 - if $pmtn \in \beta$ then processing of a job may be interrupted and resumed later even on a different machine
 - if $pmtn \notin \beta$ then the processing of an operation must not be interrupted
- $prec$: precedence constraints
 - there are precedence relations (partial order) between jobs: a job may not start before all its predecessors have been finished
 - typically represented by directed acyclic graph
 - special forms of precedence relations, e.g., *chain*, *intree*, *outtree*
- d_j : jobs have deadlines by which they must finish
→ discussion due date vs. deadline

Classification – objective function (γ)

Notation:

- C_j completion time of job j
- $F_j = C_j - r_j$ flowtime of job j
- $L_j = C_j - d_j$ lateness of job j
- $T_j = \max\{0, C_j - d_j\}$ tardiness of job j
- U_j unit penalty: if $C_j > d_j$ then $U_j = 1$, otherwise $U_j = 0$

Objective functions (min):

- C_{\max} : makespan $C_{\max} = \max_j C_j$
- F_{\max} : maximum flowtime $F_{\max} = \max_j F_j$
- L_{\max} : maximum lateness $L_{\max} = \max_j L_j$
- $\sum(w_j)C_j, \sum(w_j)F_j, \sum(w_j)T_j, \dots$
- $\sum(w_j)U_j$ (weighted) number of tardy jobs